### **SPECIAL LINEAR SYSTEMS OF EQUATIONS**

- Symmetric positive definite matrices.
- ullet The  $LDL^T$  decomposition; The Cholesky factorization
- Banded systems

6-1

### A few properties of SPD matrices

- $\triangleright$  Diagonal entries of A are positive
- Recall: the k-th principal submatrix  $A_k$  is the  $k \times k$  submatrix of A with entries  $a_{ij}$ ,  $1 \le i, j \le k$  (Matlab: A(1:k,1:k)).
- lacktriangle Each  $A_k$  is SPD
- $\triangle$  Consequence:  $Det(A_k) > 0$  for  $k = 1, \dots, n$ .
- For any  $n \times k$  matrix X of rank k, the matrix  $X^T A X$  is SPD.
- ightharpoonup The mapping :  $x,y \rightarrow (x,y)_A \equiv (Ax,y)$

defines a proper inner product on  $\mathbb{R}^n$ . The associated norm, denoted by  $\|\cdot\|_A$ , is called the energy norm, or simply the A-norm:

$$\|x\|_A = (Ax,x)^{1/2} = \sqrt{x^T A x}$$

### Positive-Definite Matrices

> A real matrix is said to be positive definite if

$$(Au,u)>0$$
 for all  $u
eq 0$   $u\in \mathbb{R}^n$ 

Let A be a real positive definite matrix. Then there is a scalar lpha>0 such that

$$(Au,u) \geq lpha \|u\|_2^2.$$

- Consider now the case of Symmetric Positive Definite (SPD) matrices.
- ➤ Consequence 1: **A** is nonsingular
- $\triangleright$  Consequence 2: the eigenvalues of A are (real) positive

6-2 TB: 23; AB:1.3.1–.2,1.5.1–4; GvL 4 – SPD

6-2

Related measure in Machine Learning, Vision, Statistics: the Mahalanobis distance between two vectors:

$$d_A(x,y) = \|x-y\|_A = \sqrt{(x-y)^T A(x-y)}$$

Appropriate distance (measured in # standard deviations) if x is a sample generated by a Gaussian distribution with covariance matrix A and center y.

# More terminology

- ➤ A matrix is Positive  $(Au,u) \geq 0$  for all  $u \in \mathbb{R}^n$ Semi-Definite if:
- Eigenvalues of symmetric positive semi-definite matrices are real nonnegative, i.e., ...
- $\triangleright$  ... A can be singular [If not, A is SPD]
- $\triangleright$  A matrix is said to be Negative Definite if -A is positive definite. Similar definition for Negative Semi-Definite
- A matrix that is neither positive semi-definite nor negative semidefinite is indefinite
- Show that if  $A^T = A$  and  $(Ax, x) = 0 \ \forall x$  then A = 0
- Show: A is indefinite iff  $\exists x,y:(Ax,x)(Ay,y)<0$

TB: 23; AB:1.3.1-.2,1.5.1-4; GvL 4 - SPD

Cholesky factorization is a specialization of the LU factorization for the SPD case. Several variants exist.

First algorithm: row-oriented LDLT

Adapted from Gaussian Elimination [Work only on upper triang. part]

- 1. For k = 1 : n 1 Do: For i = k + 1 : n Do: piv := a(k,i)/a(k,k)a(i,i:n) := a(i,i:n) - piv \* a(k,i:n)End
- This will give the U matrix of the LU factorization. Therefore  $D = diaq(U), L^T = D^{-1}U.$

# The $LDL^T$ and Cholesky factorizations

- Arr The LU factorization of an SPD matrix A exists
- $\blacktriangleright$  Let A=LU and D=diag(U) and set  $M\equiv (D^{-1}U)^T$ .

Then

$$A = LU = LD(D^{-1}U) = LDM^T$$

- $\triangleright$  Both L and M are unit lower triangular
- Consider  $L^{-1}AL^{-T} = DM^TL^{-T}$
- Matrix on the right is upper triangular. But it is also symmetric. Therefore  $M^T L^{-T} = I$  and so M = L
- $\blacktriangleright$  The diagonal entries of D are positive [Proof: consider  $L^{-1}AL^{-T}=$ D]. In the end:

$$A = LDL^T = GG^T$$
 where  $G = LD^{1/2}$ 

TB: 23; AB:1.3.1-.2,1.5.1-4; GvL 4 - SPD

Row-Cholesky (outer product form)

Scale the rows as the algorithm proceeds. Line 4 becomes

$$a(i,:) := a(i,:) - [a(k,i)/\sqrt{a(k,k)}] * \left\lceil a(k,:)/\sqrt{a(k,k)} \right\rceil$$

### ALGORITHM: 1. Outer product Cholesky

- 1. For k = 1 : n Do:
- $A(k,k:n) = A(k,k:n)/\sqrt{A(k,k)}$ ;
- For i:=k+1:n Do :
- A(i, i : n) = A(i, i : n) A(k, i) \* A(k, i : n);4.
- 5. End
- 6. End
- Result: Upper triangular matrix U such  $A = U^T U$ .

TB: 23; AB:1.3.1-.2,1.5.1-4; GvL 4 - SPD

6. End

# Example:

$$A = \begin{pmatrix} 1 & -1 & 2 \\ -1 & 5 & 0 \\ 2 & 0 & 9 \end{pmatrix}$$

- $\triangle$  Is A symmetric positive definite?
- Mhat is the  $LDL^T$  factorization of A?
- lacktriangle What is the Cholesky factorization of A?

9 TB: 23; AB:1.3.1–.2,1.5.1–4; GvL 4 – SPD

6-9

Assume that first i-1 columns of G already known.

➤ Compute unscaled column-vector:

$$v = A(:,j) - \sum_{k=1}^{j-1} G(j,k) G(:,k)$$

- ightharpoonup Notice that  $v(j)\equiv G(j,j)^2$ .
- lacksquare Compute  $\sqrt{v(j)}$  and scale v to get j-th column of G.

Column Cholesky. Let  $A = GG^T$  with G = lower triangular. Then equate j-th columns:

$$a(i,j) = \sum_{k=1}^j g(j,k) g^T(k,i) 
ightarrow$$

$$egin{align} A(:,j) &= \sum_{k=1}^{j} G(j,k) G(:,k) \ &= G(j,j) G(:,j) + \sum_{k=1}^{j-1} G(j,k) G(:,k) 
ightarrow \ G(j,j) G(:,j) &= A(:,j) - \sum_{k=1}^{j-1} G(j,k) G(:,k) \ \end{gathered}$$

6-10

TB: 23: AB:1.3.1-.2.1.5.1-4: GvL 4 - SPD

6-10

ALGORITHM: 2. Column Cholesky

- 1. For j = 1 : n do
- 2. For k = 1: j 1 do
- 3. A(j:n,j) = A(j:n,j) A(j,k) \* A(j:n,k)
- 4. EndDo
- 5. If  $A(j,j) \leq 0$  ExitError("Matrix not SPD")
- 6.  $A(j,j) = \sqrt{A(j,j)}$
- 7. A(j+1:n,j) = A(j+1:n,j)/A(j,j)
- 8. EndDo

Try algorithm on:

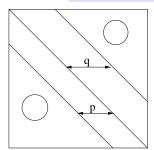
$$A = \begin{pmatrix} 1 & -1 & 2 \\ -1 & 5 & 0 \\ 2 & 0 & 9 \end{pmatrix}$$

TB: 23: AB:1.3.1-2.1.5.1-4: Gvl. 4 - SPD

TR: 23: AB:1 3 1- 2 1 5 1-4: Gvl 4 - SPD

#### Banded matrices

- ➤ Banded matrices arise in many applications
- lacksquare A has upper bandwidth q if  $a_{ij}=0$  for j-i>q
- lacksquare A has lower bandwidth p if  $|a_{ij}=0|$  for i-j>p



➤ Simplest case: tridiagonal ➤ p = q = 1.

6-13

TB: 23; AB:1.3.1-.2,1.5.1-4; GvL 4 - SPD

6-1

First observation: Gaussian elimination (no pivoting) preserves the initial banded form. Consider first step of Gaussian elimination:

2. For 
$$i = 2:n$$
 Do:  
3.  $a_{i1} := a_{i1}/a_{11}$  (pivots)  
4. For  $j := 2:n$  Do:  
5.  $a_{ij} := a_{ij} - a_{i1} * a_{1j}$   
6. End  
7. End

If A has upper bandwidth q and lower bandwidth p then so is the resulting [L/U] matrix.  $\triangleright$  Band form is preserved (induction)

Operation count?

6-14

TB: 23; AB:1.3.1-.2,1.5.1-4; GvL 4 - SPD

6-14

# What happens when partial pivoting is used?

If A has lower bandwidth p, upper bandwidth q, and if Gaussian elimination with partial pivoting is used, then the resulting U has upper bandwidth p+q. L has at most p+1 nonzero elements per column (bandedness is lost).

➤ Simplest case: tridiagonal ➤ p = q = 1.

### Example:

$$A = \begin{pmatrix} 1 & 1 & 0 & 0 & 0 \\ 2 & 1 & 1 & 0 & 0 \\ 0 & 2 & 1 & 1 & 0 \\ 0 & 0 & 2 & 1 & 1 \\ 0 & 0 & 0 & 2 & 1 \end{pmatrix}$$

TB: 23; AB:1.3.1-.2,1.5.1-4; GvL 4 - SPD