A solution that doesn’t help

- Part of what makes this classic attack easy is that the array grows in the direction toward the function’s return address
- If we made the stack grow towards higher addresses instead, this wouldn’t work in the same way
- Classic puzzler: why isn’t this a solution to the problem?

A concrete example

```c
void func(char *attacker_controlled) {
    char buffer[50];
    strcpy(buffer, attacker_controlled);
}
```

What might happen in this example, for instance?

Stack direction orientation

- Higher addresses are “deeper” in the stack, and represent older stack frames (callers) and data (pushed first)
- Lower addresses are closer to the “top” of the stack, representing more-recently pushed frames (callee) and data

Stack frame normal overflow

Reversed overflow

Outline

- Reversing the stack, discussion
- Other safety problems
- Logistics announcements
- Integer overflow example
- Code auditing
Non-contiguous overflow

- An overflow doesn't have to write to the buffer in sequence
- For instance, the code might compute a single index, and store to it

Heap buffer overflow

- Overwriting a malloc-ed buffer isn't close to a return address
- But other targets are available:
  - Metadata used to manage the heap, contents of other objects

Use after free

- A common bug is to free an object via one pointer and keep using it via another
- Leads to unsafe behavior after the memory is reused for another object

Integer overflow

- Integer types have limited size, and will wrap around if a computation is too large
- Not unsafe itself, but often triggers later bugs
  - E.g., not allocating enough space

Function pointers, etc.

- Other data used for control flow could be targeted for overwriting by an attacker
- Common C case: function pointers
- More obscure C case: set jmp/long jmp buffers

Virtual dispatch

- When C++ objects have virtual methods, which implementation is called depends on the runtime type
- Under the hood, this is implemented with a table of function pointers called a vtable
- An appealing target in attacking C++ code

Non-control data overwrite

- An attacker can also trigger undesired-to-you behavior by modifying other data
- For instance, flags that control other security checks

Format string injection

- The first argument of printf is a little language controlling output formatting
- Best practice is for the format string to be a constant
- An attacker who controls a format string can trigger other mischief
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Piazza for announcements, Q&A
- Sign up to the Piazza page for announcements, questions and answers
- Preferred over email for both announcements and Q&A
- Get to via links from Canvas or public pages
- There now: some details from Thursday's demo

Project 1 schedule
- The tentative schedule for project 1 was too optimistic
- Expect to see it by next Tuesday, other deadlines later as well

Lab participation
- In online format, err on the side of being explicit about whether we're recording your participation
- Good to check before you leave, especially if early
- Let us know if we missed your participation last week

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Integer overflow to buffer overflow
- One common pattern: overflow causes an allocation to be too small
- In machine integers, multiplication doesn't always make a value larger

Overflow example
struct obj { short ident, x, y, z; long b; double c;};
struct obj *read_objs(int num_objs) {
    unsigned int size = num_objs*(unsigned)sizeof(obj);
    struct obj *objs = malloc(size);
    struct obj *p = objs;
    for (i = 0; i < num_objs; i++) {
        fread(p, sizeof(struct obj), 1, stdin);
        if (p->ident == 0x4442) return 0;
        /* ... */ p++; }
    return objs; }

Overflow example questions
1. What's a value of num_objs that would trigger an overflow?
   - Think back to 2021 on how multiplication overflows
2. Why is the p->ident check relevant to exploitability?

http://www-users.cselabs.umn.edu/classes/Fall-2020/csci4271/slides/02/overflow-eg.c
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Auditing is...
- Reading code to find security bugs
- Threat modeling comes first, tells you what kinds of bugs you're looking for
- Bug fixing comes next (might be someone else's job)

Tiers and triage
- You might not have time to do a complete job, so use auditing time strategically
- Which bugs are most likely, and easiest to find?
- Triage into definitely safe, definitively unsafe, hard to tell
  - Hard to tell might be improved even if safe

Threat model and taint
- Vulnerability depends on what an attacker might control
- Another word for attacker-controlled is "tainted"
- Threat model is the best source of tainting information
  - Of course, can always be conservative

Where to look for problems
- If you can't read all the code carefully, search for indicators of common danger spots
  - For format strings, look for printf
  - For buffer overflows, look at buffers and copying functions

Ideal: proof
- Given enough time, for each dangerous spot, be able to convince someone:
  - Proof of safety: reasons why a bug could never happen, could turn into assertions
  - Proof of vulnerability: example of tainted input that causes a crash

Auditing exercise
- BCLPR is a buggy program from a previous year's 5271
- This code has at least three buffer overflow bugs: where are they?
- Are all the bugs exploitable? As an attacker, could you use them?

http://www-users.cselabs.umn.edu/classes/Fall-2020/csci4271/slides/02/bclpr.c