## Minimax (Ch. 5-5.3)

COMPLETE MAP OF OPTIMAL TIC-TAC-TOE MOVES
YOUR MOVE IS GIVEN BY THE POSITIN OF THE LARGEST RED SYMB $\alpha$ ON THE GRID. WHEN YOUR OPPONENT PICKS A MOVE, ZOOM IN ON THE REGION OF THE GRID WHERE THEY WENT. REPEAT.

MAP FOR X:



## Announcements

## Homework 1 solutions posted

Test next week ( $9^{\text {th }}$ )
-Covers up to and including HW2
(all search, Ch. 1-4)

## Single-agent

So far we have look at how a single agent can search the environment based on its actions

Now we will extend this to cases where you are not the only one changing the state (i.e. multi-agent)

The first thing we have to do is figure out how to represent these types of problems

## Multi-agent (competitive)

Most games only have a utility (or value) associated with the end of the game (leaf node)

So instead of having a "goal" state (with possibly infinite actions), we will assume:
(1) All actions eventually lead to terminal state (i.e. a leaf in the tree)
(2) We know the value (utility) only at leaves

## Multi-agent (competitive)

For now we will focus on zero-sum two-player games, which means a loss for one person is a gain for another

Betting is a good example of this: If I win I get $\$ 5$ (from you), if you win you get $\$ 1$ (from me). My gain corresponds to your loss

Zero-sum does not technically need to add to zero, just that the sum of scores is constant

## Multi-agent (competitive)

Zero sum games mean rather than representing outcomes as:
[ $\mathrm{Me}=5$, You $=-5$ ]
We can represent it with a single number: [ $\mathrm{Me}=5$ ], as we know: $\mathrm{Me}+\mathrm{You}=0$ (or some c)

This lets us write a single outcome which "Me" wants to maximize and "You" wants to minimize

## Minimax

Thus the root (our agent) will start with a maximizing node, the the opponent will get minimizing noes, then back to max... repeat...

This alternation of maximums and minimums is called minimax

I will use $\triangle$ to denote nodes that try to maximize and $\nabla$ for minimizing nodes

## Minimax

Let's say you are treating a friend to lunch. You choose either: Shuang Cheng or Afro Deli

The friend always orders the most inexpensive item, you want to treat your friend to best food

Which restaurant should you go to?
Menus:
Shuang Cheng: Fried Rice=\$10.25, Lo Mein=\$8.55 Afro Deli: Cheeseburger=\$6.25, Wrap=\$8.74

## Minimax

## Shuang Cheng

## Lo Mein

8.55

### 10.25

## Cheese- burger

## Minimax

You could phrase this problem as a set of maximum and minimums as: $\max (\min (8.55,10.25), \min (6.25,8.55))$
... which corresponds to: max( Shuang Cheng choice, Afro Deli choice)

If our goal is to spend the most money on our friend, we should go to Shuang Cheng

## Minimax

One way to solve this is from the leaves up:


## Minimax

$\max (\min (1,3), 2, \min (0,4))=2$, should pick Order: $\quad$ action F $1^{\text {st. }}$ R (can swap $2^{2^{\text {nd }} .} . \mathrm{B} B$ and R )

## Minimax



## Minimax

This representation works, but even in small games you can get a very large search tree

For example, tic-tac-toe has about 9! actions to search (or about 300,000 nodes)

Larger problems (like chess or go) are not feasible for this approach (more on this next class)

## Minimax

## "Pruning" in real life:

## Snip branch


"Pruning" in CSCI trees:

Snip branch


## Alpha-beta pruning

However, we can get the same answer with searching less by using efficient "pruning"

It is possible to prune a minimax search that will never "accidentally" prune the optimal solution

A popular technique for doing this is called alpha-beta pruning (see next slide)

## Alpha-beta pruning

Consider if we were finding the following: $\max (5, \min (3,19))$

There is a "short circuit evaluation" for this, namely the value of 19 does not matter
$\min (3, x) \leq 3$ for all $x$
Thus $\max (5, \min (3, x))=5$ for any $x$
Alpha-beta pruning would not search x above

## Alpha-beta pruning

If when checking a min-node, we ever find a value less than the parent's "best" value, we can stop searching this branch


## Alpha-beta pruning

In the previous slide, "best" is the "alpha" in the alpha-beta pruning
(Similarly the "worst" in a min-node is "beta")
Alpha-beta pruning algorithm:
Do minimax as normal, except: min node: if parent's "best" value greater than current node, stop \& tell parent current value max node: if parent's "worst" value less than current node, stop search and return current

## Alpha-beta pruning

## Let's solve this with alpha-beta pruning



## Alpha-beta pruning

$\max (\min (1,3), 2, \min (0, ? ?))=2$, should pick Order:
$1^{\text {st. }}$. Red $2^{\text {nd }}$. Blue $3^{\text {rd }}$. Purp
 action F

Do not
consider

## Alpha-beta pruning

## \rantOn

## I think the book is

 confusing aboutalpha-beta, especially Figure 5.5


## $\alpha \beta$ pruning



## Alpha-beta pruning

In general, alpha-beta pruning allows you to search to a depth 2d for the minimax search cost of depth d

So if minimax needs to find: $O\left(b^{m}\right)$ Then, alpha-beta searches: $\mathrm{O}\left(\mathrm{b}^{\mathrm{m} / 2}\right)$

This is exponentially better, but the worst case is the same as minimax

## Alpha-beta pruning

Ideally you would want to put your best (largest for max, smallest for min) actions first

This way you can prune more of the tree as a min node stops more often for larger "best"

Obviously you do not know the best move, (otherwise why are you searching?) but some effort into guessing goes a long way
(i.e. exponentially less states)

## Side note:

In alpha-beta pruning, the heuristic for guess which move is best can be complex, as you can greatly effect pruning

While for A* search, the heuristic had to be very fast to be useful
(otherwise computing the heuristic would take longer than the original search)

