Solution:

We will show that Gram-Schmidt breaks down iff the x_i 's are linearly dependent.

The only way in which GS can break down is if $r_{jj} = \|\hat{q}\|_2$ in line 7 is zero.

The main observation is that the vector \hat{q} at the end of the loop starting in line 4 (ending in line 6) is a linear combination of the x_i 's for $i = 1, \dots, j - 1$. [simple proof by induction – omitted.]

∞3 Cost of Gram-Schmidt?

Solution: Step j of the algorithm costs : $(j-1) \times 2m$ operations for line 3, $+(j-1) \times 2m$ operations for loop in line 4 + 3m operations in Lines 7 and 8 together. Total for step j = $c_j = (4j-1)m$. Total over the n columns $= T(n) = (2n^2+n)m pprox 2n^2m$.

Note: this is linear in m (number of rows) and quadratic in n (number of columns).

What is the cost of solving a linear system with the QR factorization?

Solution: According to the previous question we have a cost of $2n^3$ for the factorization (since m = n), to which we need to add the cost of solving a triangular solve $O(n^2)$ and the cost for computing $Q^T b$ which is again $O(n^2)$. In the end the cost is dominated by the QR factorization which is $2n^3$. This is 3 times more expensive than GE.