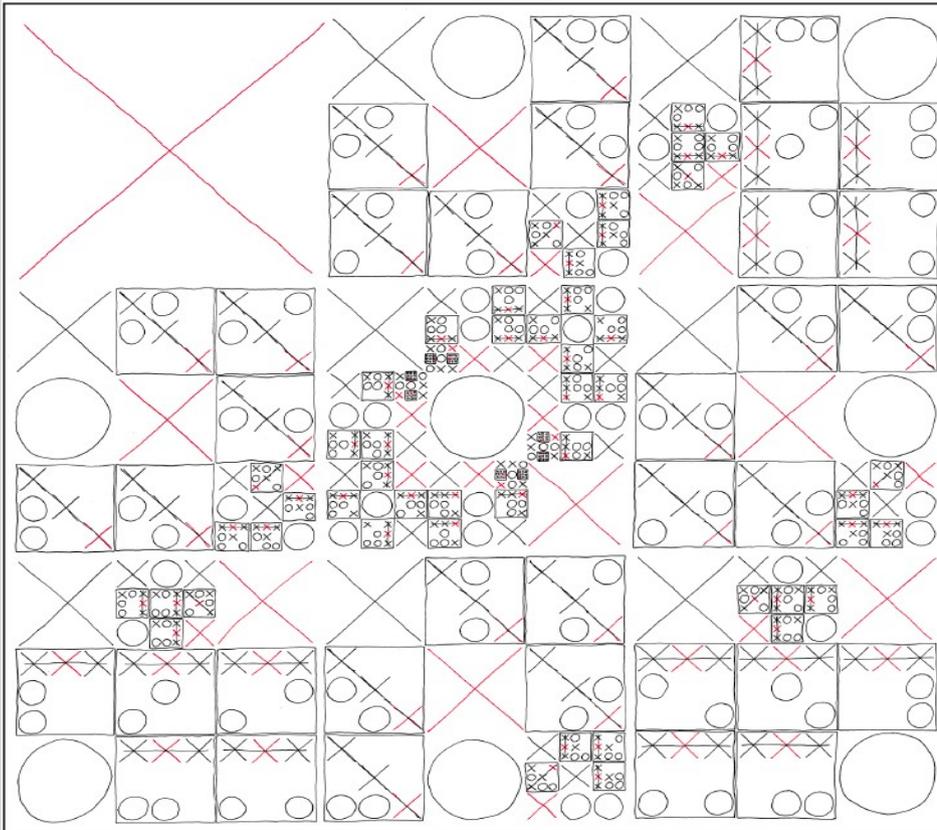


Minimax (Ch. 5-5.3)

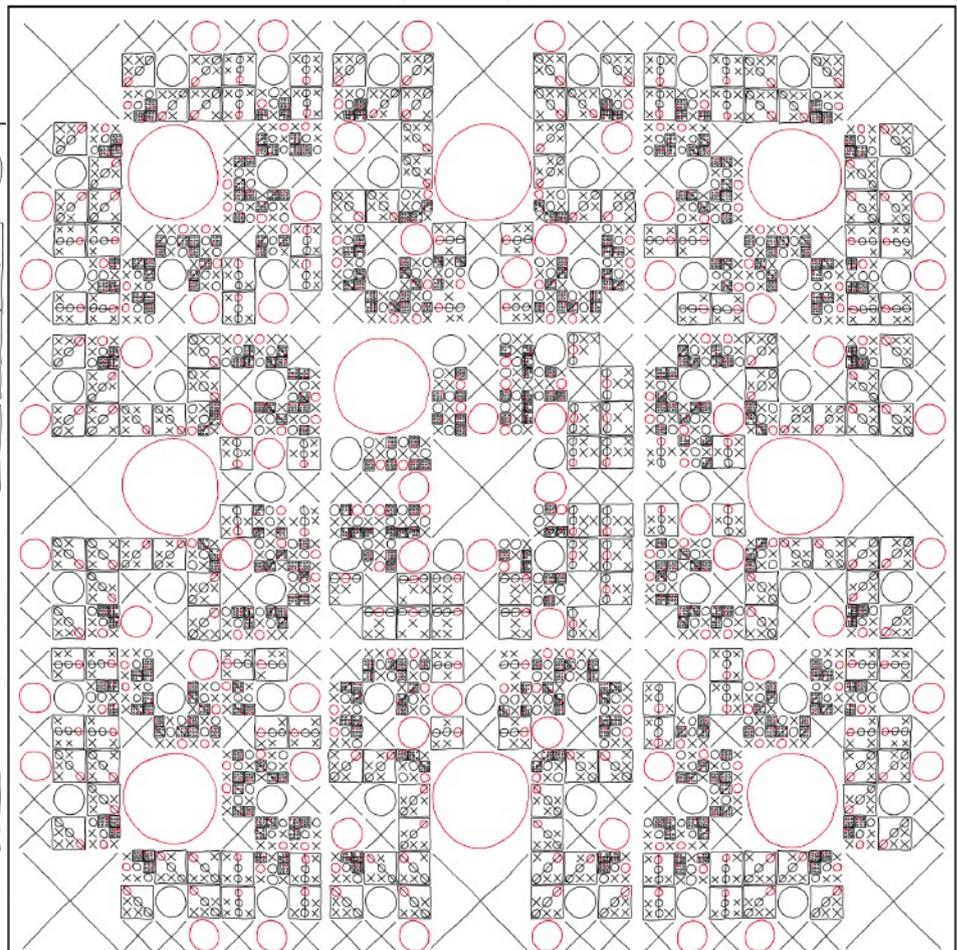
COMPLETE MAP OF OPTIMAL TIC-TAC-TOE MOVES

YOUR MOVE IS GIVEN BY THE POSITION OF THE LARGEST RED SYMBOL ON THE GRID. WHEN YOUR OPPONENT PICKS A MOVE, ZOOM IN ON THE REGION OF THE GRID WHERE THEY WENT. REPEAT.

MAP FOR X:



MAP FOR O:



Announcements

Test next week (March 4th)

-Covers up to and including HW2
(all search, Ch. 1-4)

Single-agent

So far we have look at how a single agent can search the environment based on its actions

Now we will extend this to cases where you are not the only one changing the state (i.e. multi-agent)

The first thing we have to do is figure out how to represent these types of problems

Multi-agent (competitive)

Most games only have a utility (or value) associated with the end of the game (leaf node)

So instead of having a “goal” state (with possibly infinite actions), we will assume:

(1) All actions eventually lead to terminal state (i.e. a leaf in the tree)

(2) We know the value (utility) only at leaves

Multi-agent (competitive)

For now we will focus on zero-sum two-player games, which means a loss for one person is a gain for another

Betting is a good example of this: If I win I get \$5 (from you), if you win you get \$1 (from me). My gain corresponds to your loss

Zero-sum does not technically need to add to zero, just that the sum of scores is constant

Multi-agent (competitive)

Zero sum games mean rather than representing outcomes as:

[Me=5, You =-5]

We can represent it with a single number:

[Me=5], as we know: $Me + You = 0$ (or some c)

This lets us write a single outcome which “Me” wants to maximize and “You” wants to minimize

Minimax

Thus the root (our agent) will start with a maximizing node, then the opponent will get minimizing nodes, then back to max... repeat...

This alternation of maximums and minimums is called minimax

I will use \triangle to denote nodes that try to maximize and ∇ for minimizing nodes

Minimax

Let's say you are treating a friend to lunch.
You choose either: Shuang Cheng or Afro Deli

The friend always orders the most inexpensive item, you want to treat your friend to best food

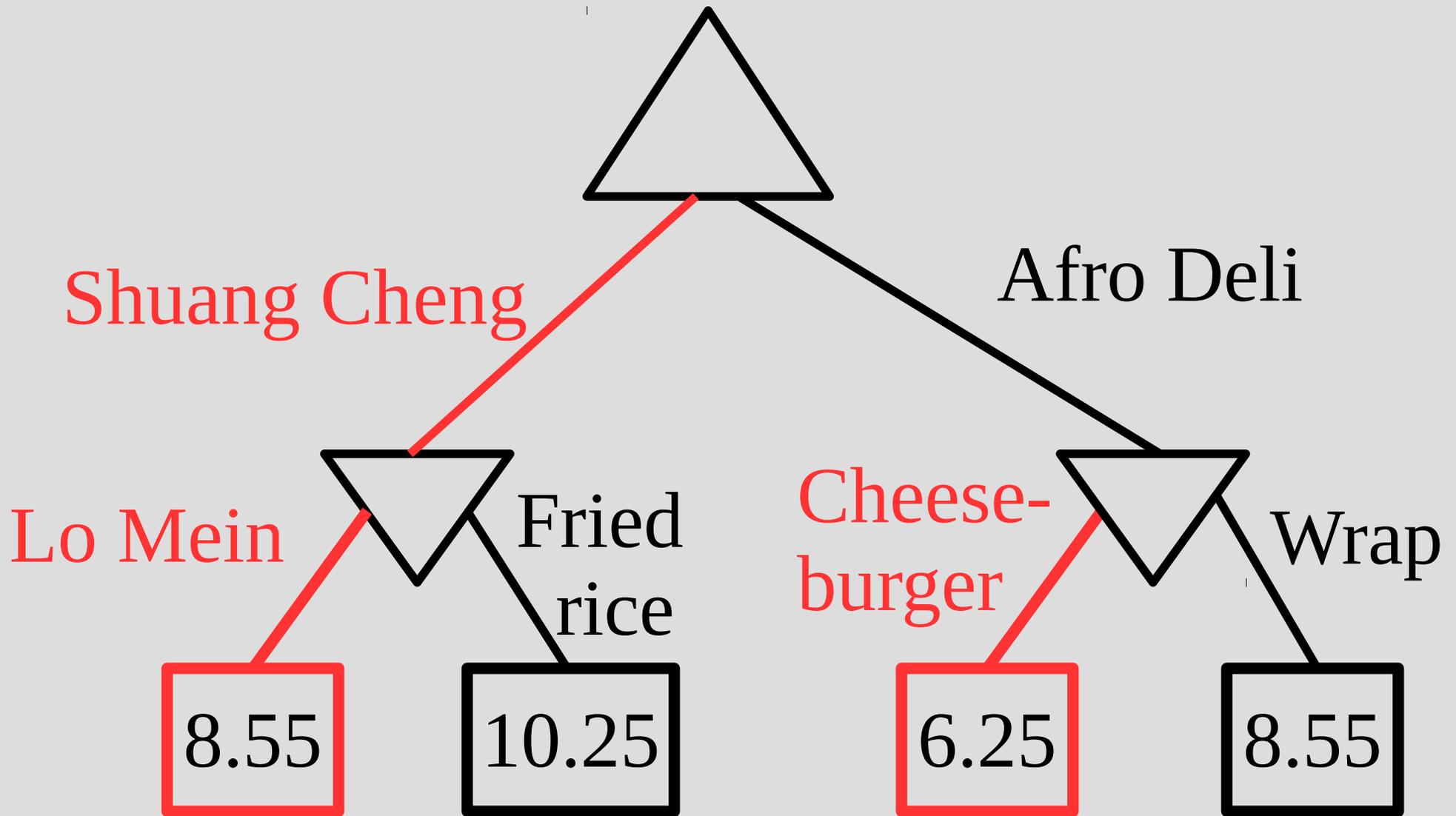
Which restaurant should you go to?

Menus:

Shuang Cheng: Fried Rice=\$10.25, Lo Mein=\$8.55

Afro Deli: Cheeseburger=\$6.25, Wrap=\$8.74

Minimax



Minimax

You could phrase this problem as a set of maximum and minimums as:

$\max(\min(8.55, 10.25), \min(6.25, 8.55))$

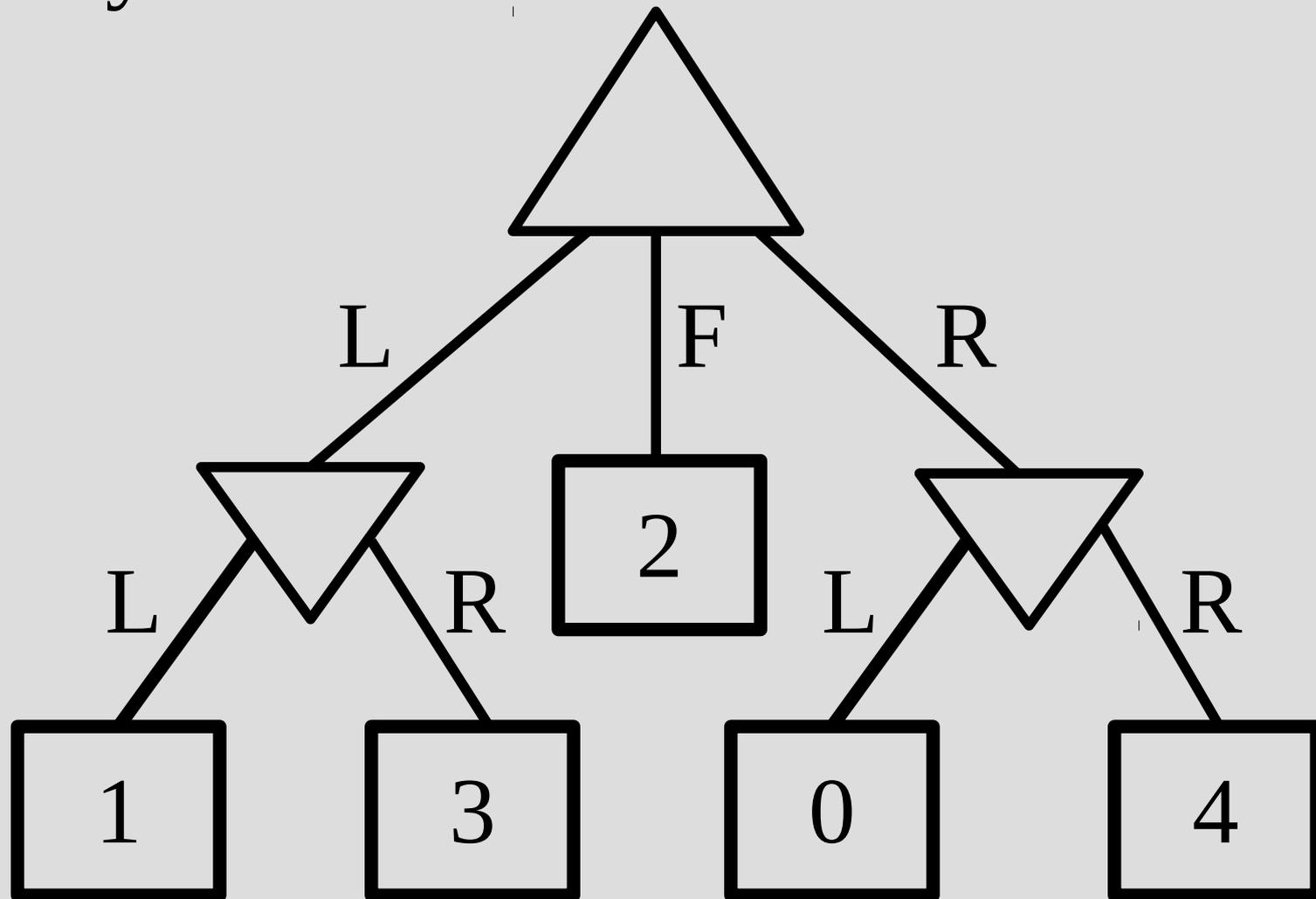
... which corresponds to:

$\max(\text{Shuang Cheng choice}, \text{Afro Deli choice})$

If our goal is to spend the most money on our friend, we should go to Shuang Cheng

Minimax

One way to solve this is from the leaves up:



Minimax

$\max(\min(1,3), 2, \min(0, 4)) = 2$, should pick

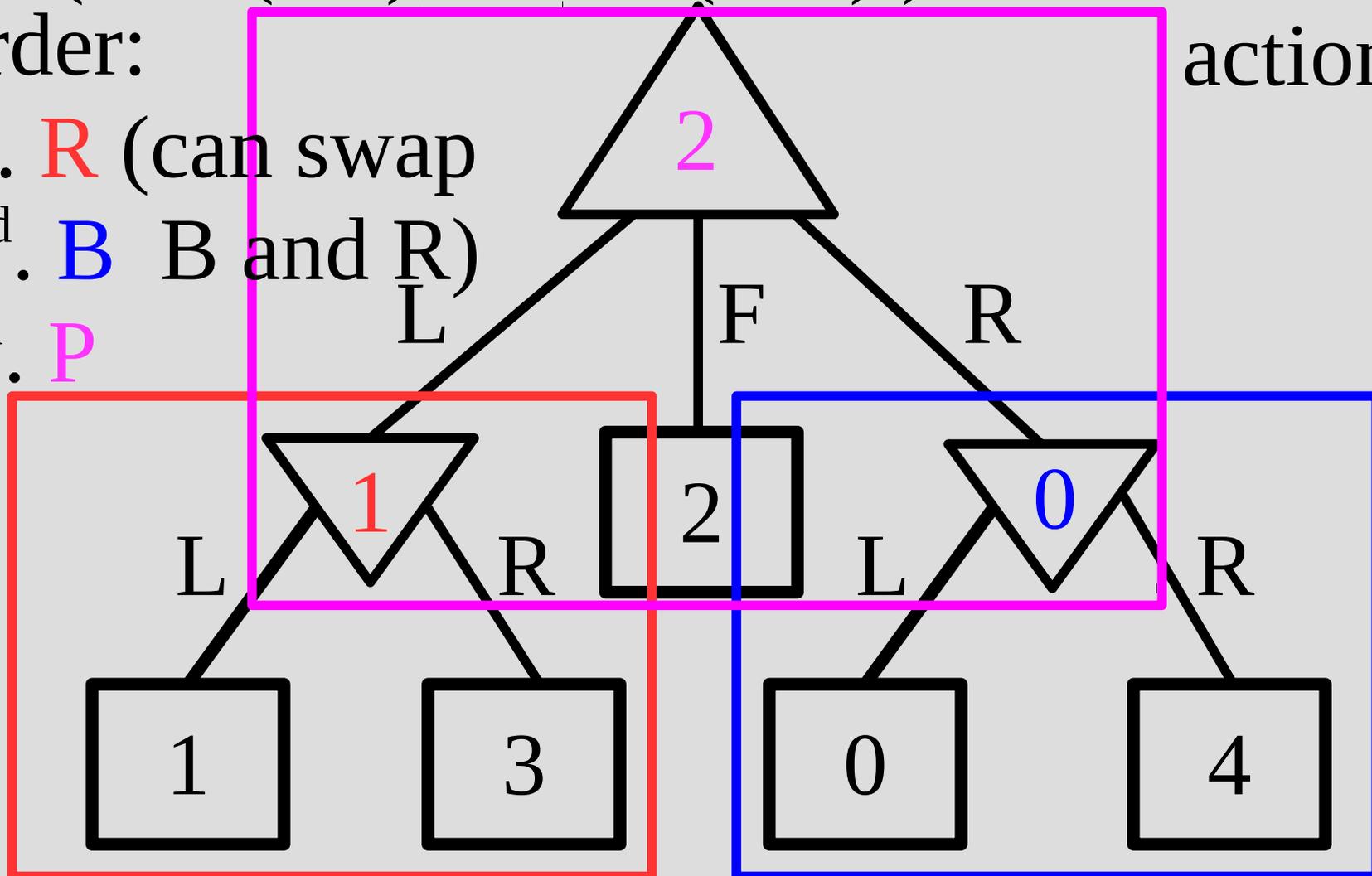
Order:

1st. **R** (can swap

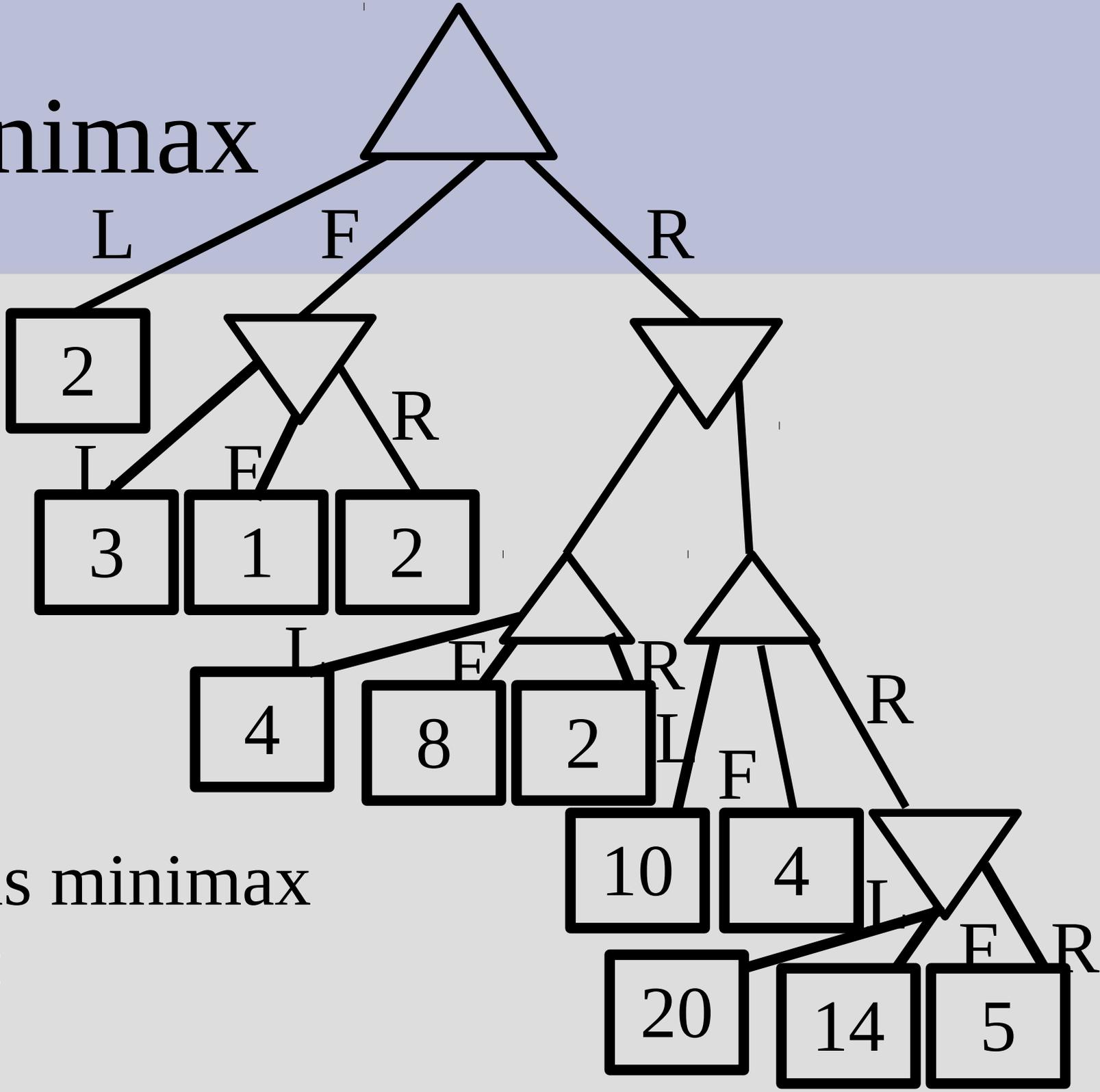
2nd. **B** B and R)

3rd. **P**

action F



Minimax



Solve this minimax problem:

Minimax

This representation works, but even in small games you can get a very large search tree

For example, tic-tac-toe has about $9!$ actions to search (or about 300,000 nodes)

Larger problems (like chess or go) are not feasible for this approach (more on this next class)

Minimax

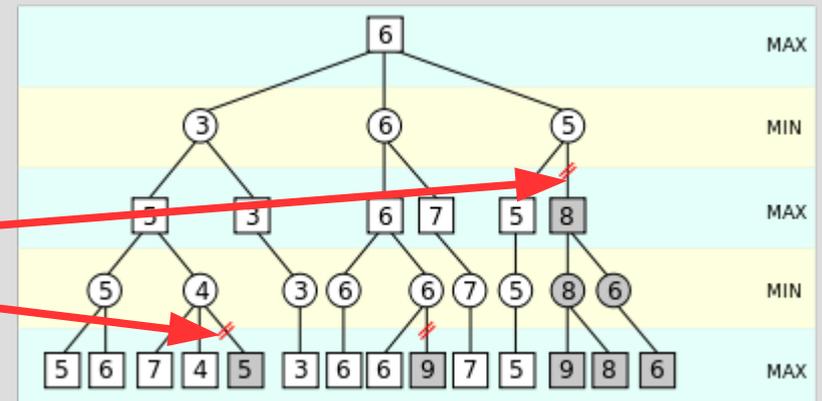
“Pruning” in real life:

Snip branch



“Pruning” in CSCCI trees:

Snip branch



Alpha-beta pruning

However, we can get the same answer with searching less by using efficient “pruning”

It is possible to prune a minimax search that will never “accidentally” prune the optimal solution

A popular technique for doing this is called alpha-beta pruning (see next slide)

Alpha-beta pruning

Consider if we were finding the following:
 $\max(5, \min(3, 19))$

There is a “short circuit evaluation” for this, namely the value of 19 does not matter

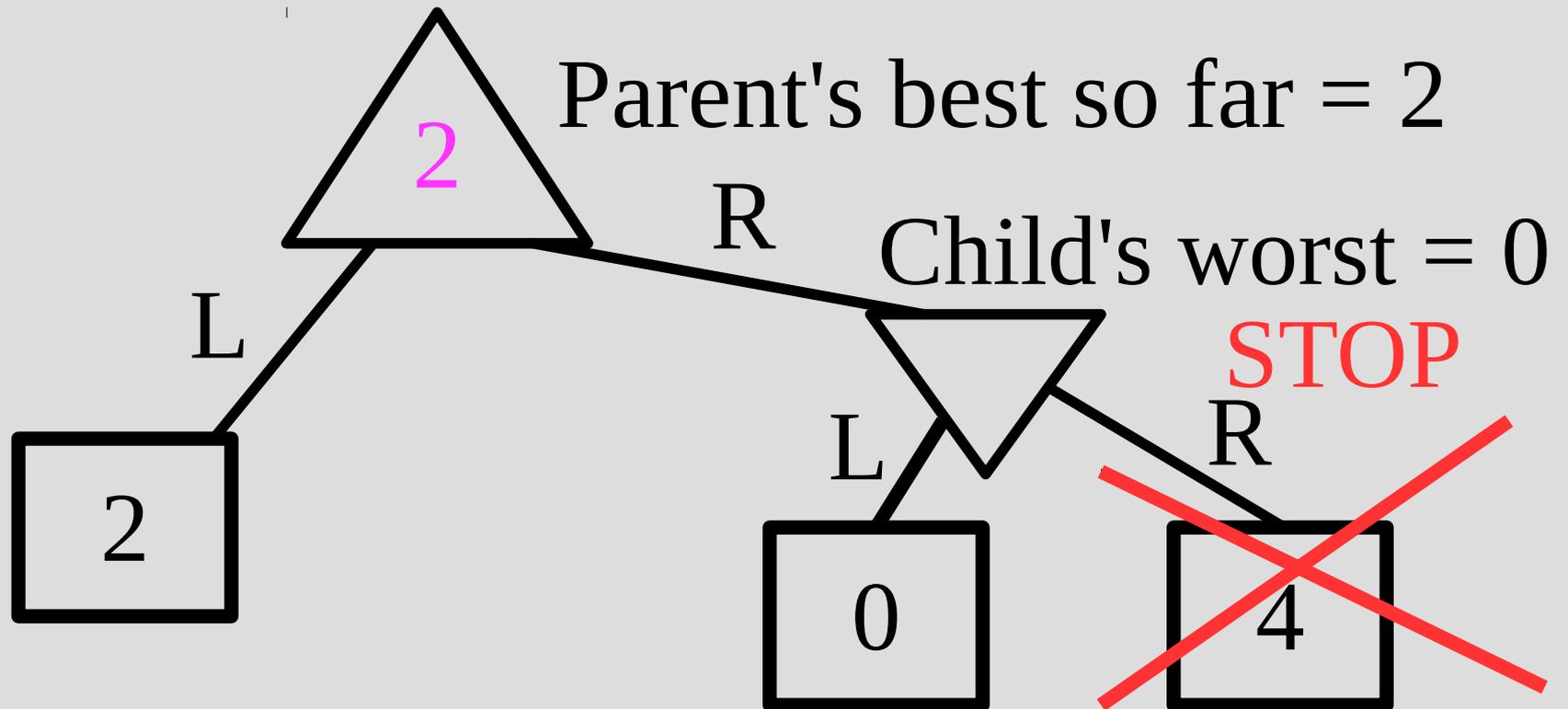
$\min(3, x) \leq 3$ for all x

Thus $\max(5, \min(3, x)) = 5$ for any x

Alpha-beta pruning would not search x above

Alpha-beta pruning

If when checking a min-node, we ever find a value less than the parent's "best" value, we can stop searching this branch



Alpha-beta pruning

In the previous slide, “best” is the “alpha” in the alpha-beta pruning (Similarly the “worst” in a min-node is “beta”)

Alpha-beta pruning algorithm:

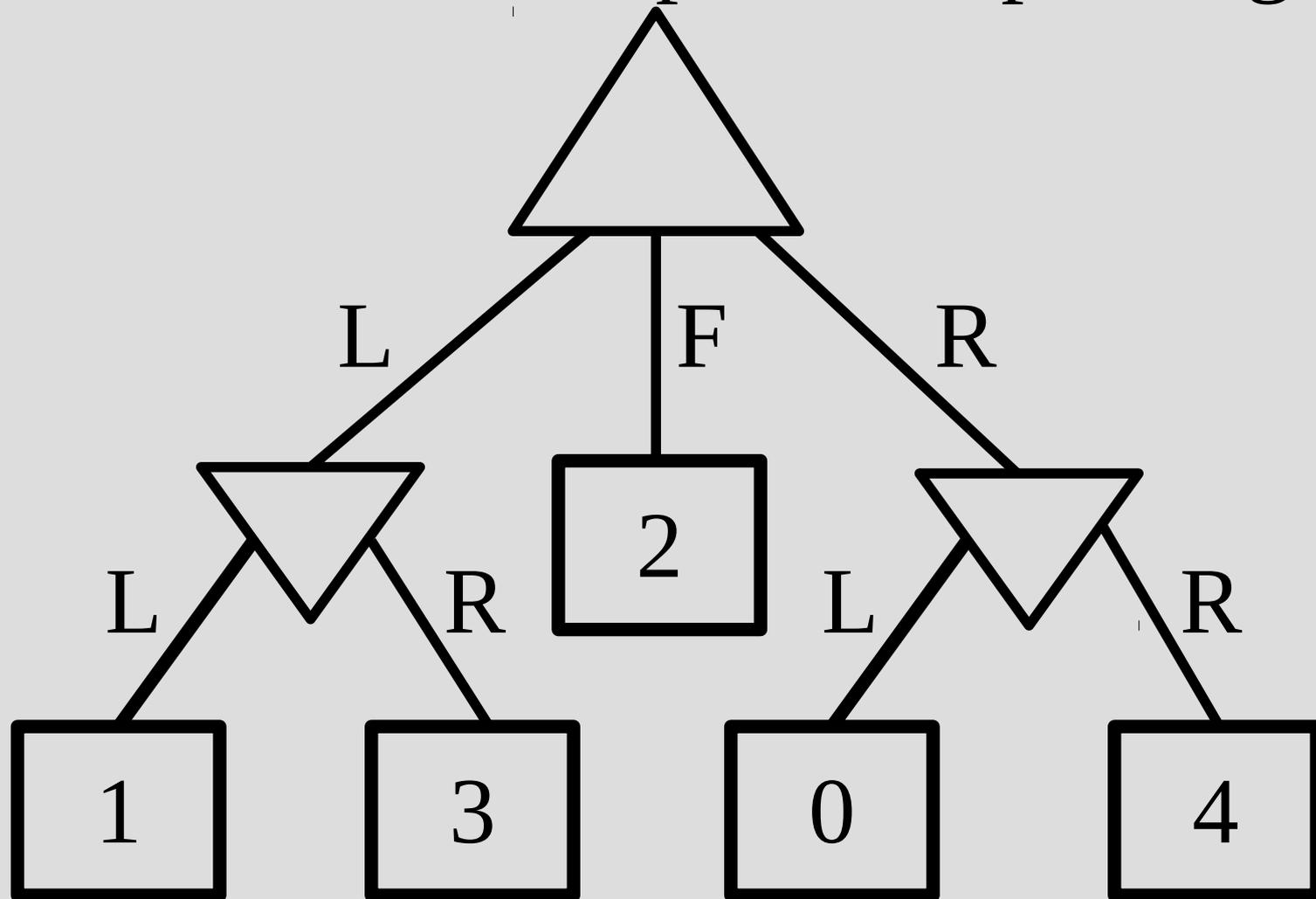
Do minimax as normal, except:

min node: if parent's “best” value greater than current node, stop & tell parent current value

max node: if parent's “worst” value less than current node, stop search and return current

Alpha-beta pruning

Let's solve this with alpha-beta pruning



Alpha-beta pruning

$\max(\min(1,3), 2, \min(0, ??)) = 2$, should pick action F

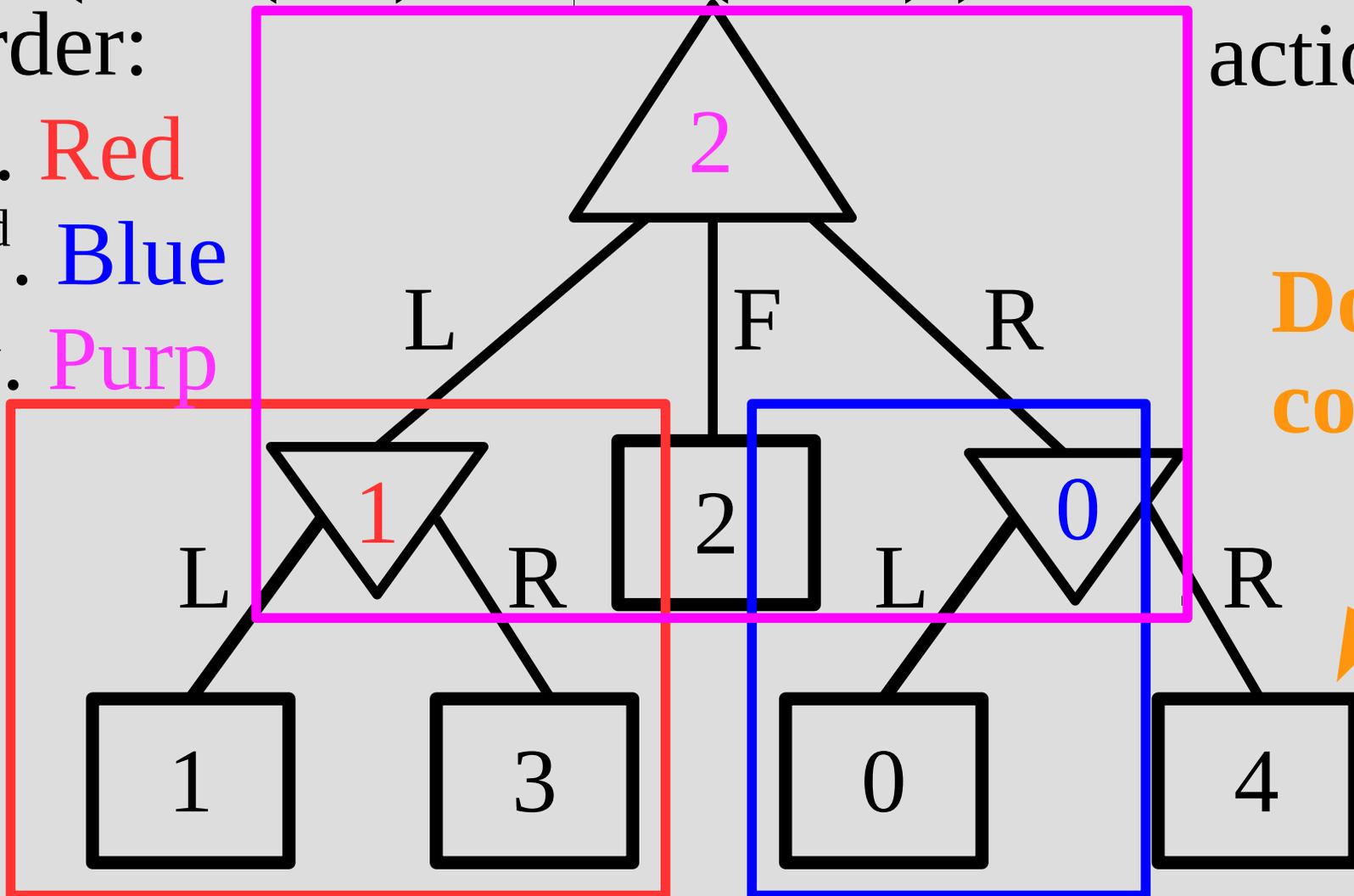
Order:

1st. Red

2nd. Blue

3rd. Purp

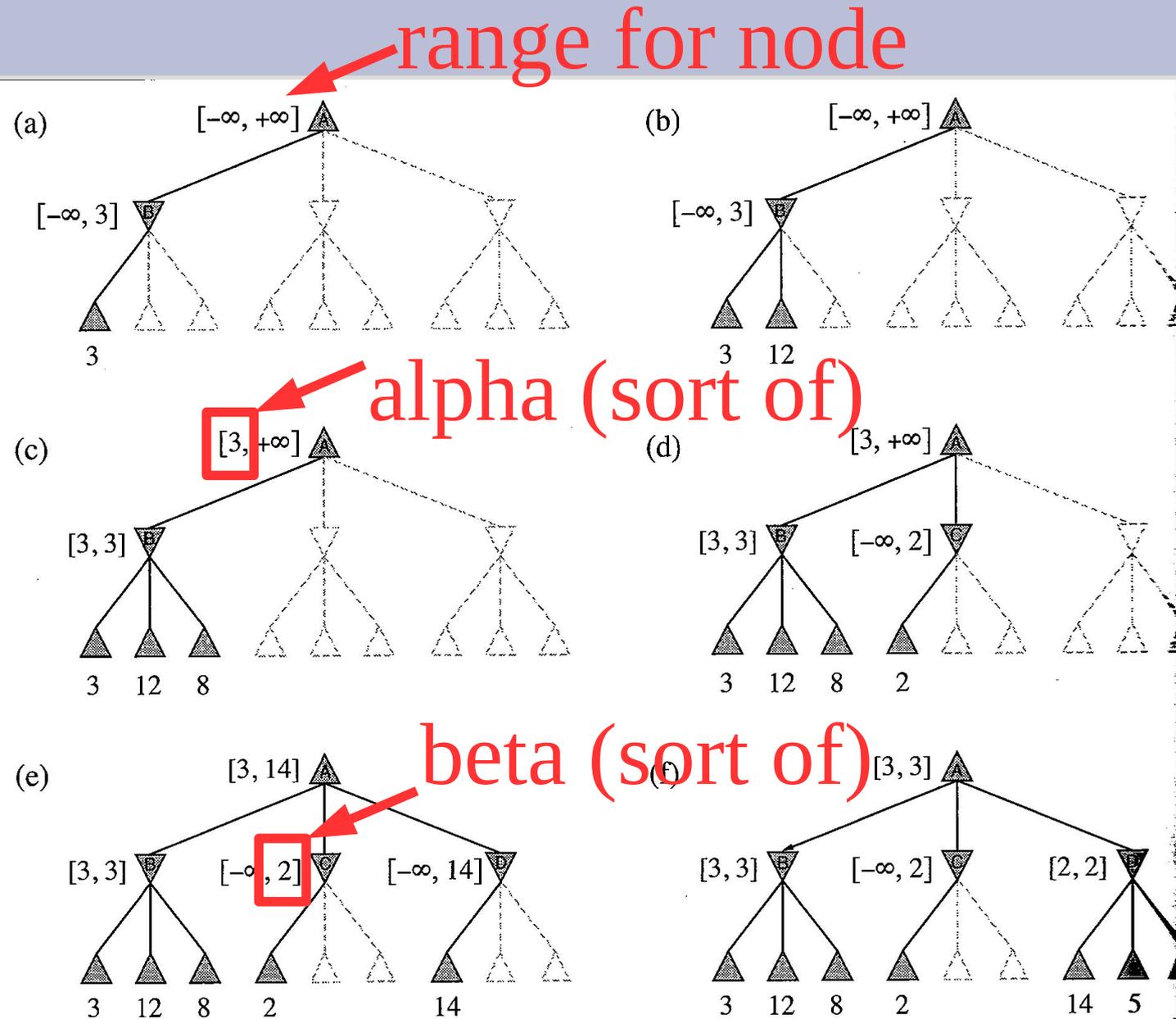
action F



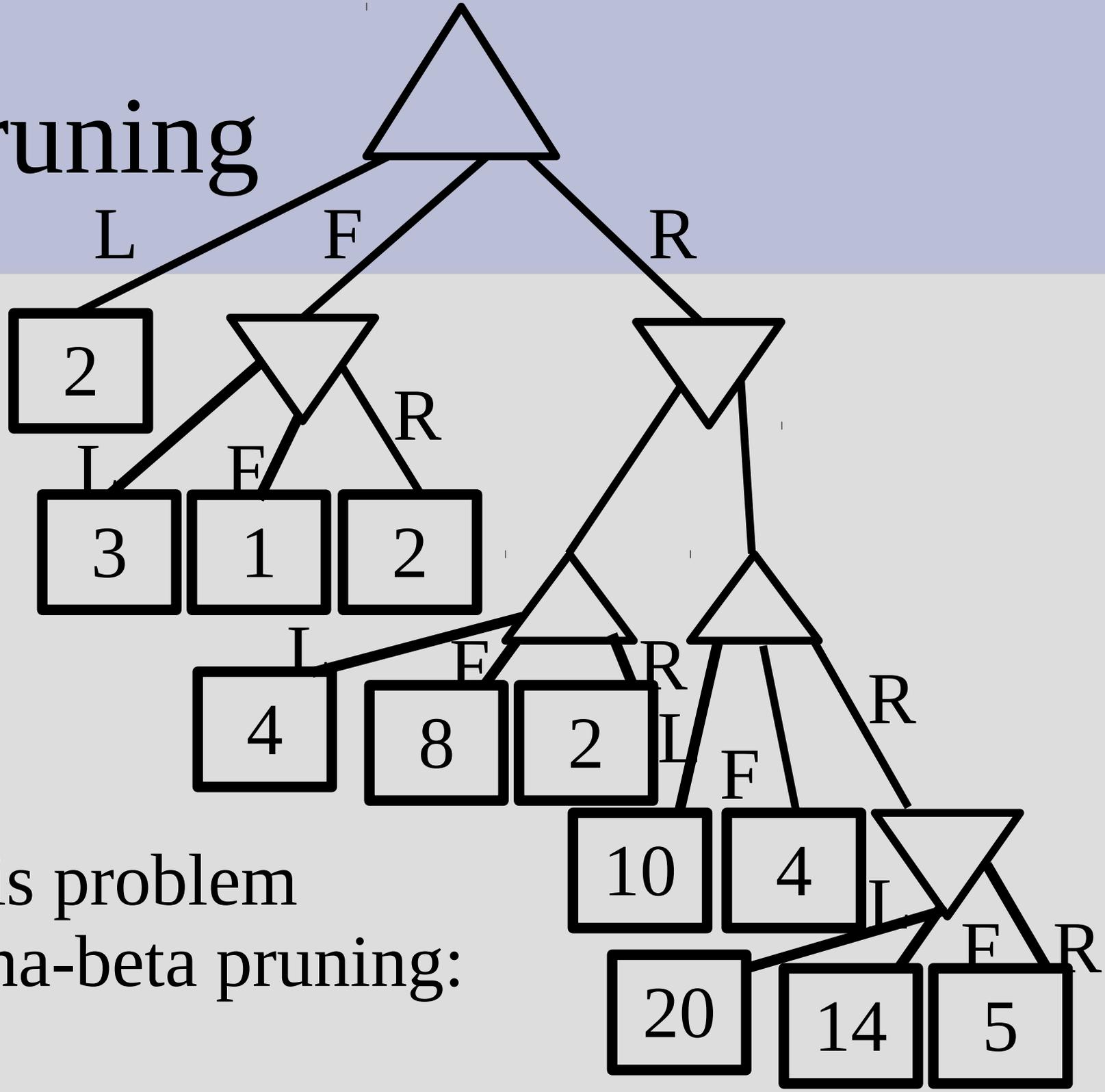
Alpha-beta pruning

\rantOn

I think the book is confusing about alpha-beta, especially Figure 5.5



$\alpha\beta$ pruning



Solve this problem
with alpha-beta pruning:

Alpha-beta pruning

In general, alpha-beta pruning allows you to search to a depth $2d$ for the minimax search cost of depth d

So if minimax needs to find: $O(b^m)$

Then, alpha-beta searches: $O(b^{m/2})$

This is exponentially better, but the worst case is the same as minimax

Alpha-beta pruning

Ideally you would want to put your best (largest for max, smallest for min) actions first

This way you can prune more of the tree as a min node stops more often for larger “best”

Obviously you do not know the best move, (otherwise why are you searching?) but some effort into guessing goes a long way (i.e. exponentially less states)

Side note:

In alpha-beta pruning, the heuristic for guess which move is best can be complex, as you can greatly effect pruning

While for A^* search, the heuristic had to be very fast to be useful
(otherwise computing the heuristic would take longer than the original search)