

## CSci Spring 2020 Section 010 Problem Set 2

Due in plain text or PDF format on Canvas at the beginning of lecture (3:35pm) on Wednesday, April 8th, 2020. We recommend that you type your solutions with a text editor or word processor and then convert them to PDF. Please label your assignment with your name, UMN email address, and the time of your lab section (12:10 pm, 1:25 pm, or 2:30 pm).

### Problem 1

**Background on division:** This problem uses the `idivq` instruction for division that has not been covered in lecture. You can read more about it in Section 3.5.5 of the textbook, and we'll describe the basics here. Here is a brief description of how 64-bit signed division works in x86-64:

- The `idivq` instruction can actually support a 128-bit dividend, with the high 64 bits in `%rdx` and the low 64 bits in `%rax`.
- A 64-bit dividend should be sign-extended to make the right 128-bit dividend.
- Start with the 64-bit dividend in `%rax`.
- Sign-extend from `%rax` to `%rdx` using the `cqto` instruction.
- The divisor can be placed in any 64-bit register.
- `idivq` takes one register as an argument which is the register that has the divisor in it.
- After `idivq`, the `%rax` register holds the quotient, while the `%rdx` register holds the remainder.

Fill in the blanks of the assembly code generated from the following C function and explain what the function does, what the parameters and variables are, and what conclusions can be made based on the return value of the function. Assume 64-bit operations and the first argument register in the assembly code contains `long n` at the start of the program. In answer, use the letters a through i to label the values you fill in the blanks. For the short answer part j, be precise.

```
int function_c(long n){
    if(n<=1){
        return 0;
    }
    long i=2;
    while(i<=n/2){
        if(n%i==0)
            return 0;
        else
            i++;
    }
    return 1;
}

function_asm:
    cmpq    $1, %rdi
    ---    .E2          a.
    movq    %rdi, %rax
    cqto
    movq    $2, %rsi
    idivq   %rsi
    ---    %rax, %r11   b.
.L1:
    cmpq    %r11, %rsi
    ---    .E1          c.
    movq    %rdi, %rax
    cqto
    idivq   %rsi
    ---    $0, %rdx    d.
    je     .E2
    addq    ----, %rsi  e.
    jmp    .L1

.E1:
    movq    ----, %rax  f.
    ----
    g.
.E2:
    ----    $0, %rax    h.
    ----
    i.
```

j. What does this function do?

## Problem 2

Consider the table below, which shows the initial contents of some registers and memory locations:

Initial Values			
Registers	Values	Memory	Values
rax	16	0x3FF0	10
rdx	32	0x3FF8	100
rcx	2	0x4000	210
rbx	0x3FF8	0x4008	24

a. Fill in Table 1 showing the results if the following machine code is run from the initial state (these instructions do not change the memory):

```
movq $1, %rax
subq $24, %rdx
addq %rcx, %rax
shlq $3, %rdx
```

Table 1			
Registers	Values	Memory	Values
rax		0x3FF0	10
rdx		0x3FF8	100
rcx		0x4000	210
rbx		0x4008	24

b. Fill in Table 2 showing the results if instead the following machine code is run from the initial state at the top:

```
leaq (%rbx, %rcx, 4), %rax
movq %rdx, 8(%rax)
subq $8, %rbx
subq $10, (%rbx)
subq $16, %rax
```

Table 2			
Registers	Values	Memory	Values
rax		0x3FF0	
rdx		0x3FF8	
rcx		0x4000	
rbx		0x4008	

### Problem 3

This is the assembly associated with the function `long function_A(long n)`:

```
function_A:
    movq    $-1, %rax
    movq    $0, %rcx
    movq    $3, %r10
    cmpq    %rcx, %rdi
    jl     .L5
    movq    $1, %rax
    movq    $1, %rdx
    jmp     .L3
.L4:
    imulq   %r10, %rax
    addq    $1, %rdx
.L3:
    cmpq    %rdi, %rdx
    jle    .L4
.L5:
    ret
```

- A. Write C code that corresponds to the assembly given above. Give the variables meaningful names, not the names of registers, including giving a more informative name for the parameter currently named `n`.
- B. Explain in a sentence or two what this function does.

## Problem 4

(Based on the textbook problem 2.87.)

We've defined a new floating point standard, called UMN-20, which follows the basic rules of IEEE floating point, but contains 20 bits. This format has 1 sign bit, 6 exponent bits ( $k=6$ ), and 13 fraction bits ( $n=13$ ). The exponent bias is  $2^{6-1} - 1 = 31$ .

A. Fill in the table that follows for each of the numbers given, with the following instructions for each column:

- **Hex**: the five hexadecimal digits describing the encoded form.
- **M**: the value of the significand. This should be a number of the form  $x$  or  $x/y$  where  $x$  is an integer and  $y$  is an integral power of 2. Examples include 1,  $67/64$ , and  $3/2$ .
- **E**: the integer value of the exponent.
- **V**: the numeric value represented. Use the notation  $\pm x$  or  $\pm x \times 2^z$ , where  $x$  and  $z$  are integers.
- **D**: the (possibly approximate) decimal numeric value. Include at least 3 non-zero fraction digits.

You need not fill in entries marked  $-$ .

**Example:** to represent the number  $3/4$  we would have  $s=0$ ,  $M=3/2$ , and  $E=-1$ . Our number would therefore have an exponent field of  $011110_2$  (decimal value of  $-1 + 31 = 30$ ) and a significand field of  $1000000000000_2$ , giving a hex representation  $3D000$ . The numeric value is  $0.75$ .

Description	Hex	M	E	V	D
$3/4$	3D000	$3/2$	-1	$3 \times 2^{-2}$	0.75
100					
Largest value $< -2$					
Smallest positive normalized value					
Number with hex 12340	12340				
NaN		-	-	-	-

B. Floating point numbers in general, and in this case specifically the UMN-20 format, support addition and subtraction, but this operation is not necessarily associative. In real numbers,  $(2^{31} + 2^{31}) - 2^{31} = 2^{31} + (2^{31} - 2^{31})$ , but suppose we did the same operations with UMN-20 floating point. Fill in the Value parts of the following table with how those results would be represented in UMN-20, briefly describing how you get each result.

Computation	Value	Computation	Value
$L_1 = 2^{31} + 2^{31}$		$R_1 = 2^{31}$	
$L_2 = 2^{31}$		$R_2 = 2^{31} - 2^{31}$	
$L = L_1 - L_2$		$R = R_1 + R_2$	

Do  $L$  and  $R$  have the same value?

## Problem 5

The assembly for the following function was produced with GCC.

```
    pushq   %rbp
    movq   %rsp, %rbp
    subq   $32, %rsp
    movq   %rdi, -24(%rbp)
    movl   $0, -8(%rbp)
.L6:
    cmpl   $25, -8(%rbp)
    jg     .L7
    movl   $26, -4(%rbp)
.L5:
    cmpl   $25, -4(%rbp)
    jle   .L4
    call   rand
    movl   %eax, -4(%rbp)
    andl   $31, -4(%rbp)
    jmp    .L5
.L4:
    movq   -24(%rbp), %rdx
    movl   -4(%rbp), %ecx
    movl   -8(%rbp), %eax
    movl   %ecx, %esi
    movl   %eax, %edi
    call   swap
    addl   $1, -8(%rbp)
    jmp    .L6
.L7:
    nop
    leave
    ret
```

Fill in the blanks for the C code, which was compiled to obtain this function. Assume that the function `rand` (part of the standard library) returns a random non-negative `int`. The function `swap(x, y, arr)` switches the position of two entries at indexes `x` and `y` in the array `arr`. You may find it helpful to make a table showing which stack locations are used to hold various local variables.

```
void create_shuffle(char *table){
    for (int i = ___; i < ___; i++){
        int j = ___;
        while (j >= ___){
            j = ___;
            j = ___ & ___;
        }
        swap(___, ___, table);
    }
}
```