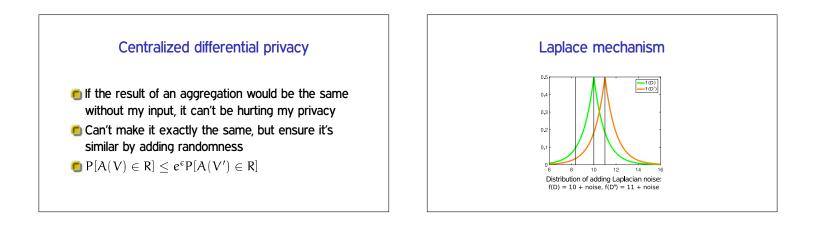
CSci 8271 Security and Privacy in Computing Day 9: Local Differential Privacy with RAPPOR Stephen McCamant

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Protecting database privacy

- Can we protect sensitive collected information to allow statistics
 - By transforming data, and/or restricting queries
- A weakness of many early attempts was breaking depending on extra information an attacker might have





 ${\color{black} \bullet}$ The parameter ${\color{black} \bullet}$ represents a privacy budget

- Often little specific guidance on choosing it
- In an interactive system, it can run out

(e, δ) differential privacy also allows a possibility of complete failure

Local differential privacy

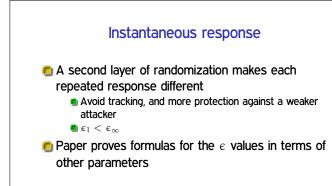
- If no trusted third party, data owners must each add their own noise
- Allows more applications, but has a worse privacy/utility tradeoff

Randomized response intuition

- Earlier proposed for embarrassing survey questions
- Randomly choose to answer either randomly or honestly
- The effect of the random answers can be removed after aggregation
- But no one can tell for sure about any particular response

Permanent response

- Repeatedly adding different noise to the same honest value would give it away
- So, add one level of noise permanently, and save the result
- Still not enough to protect "what is your age in days today?"



Some empirical results

- \blacksquare N responses let you learn at most $\sqrt{N}/10$ most common values
- In a sample distribution, detects mostly the most common elements
- Short case studies of malware binaries on Windows and Chrome user home pages