CSci 4271W
Development of Secure Software Systems
Day 9: More defenses, fuzzing

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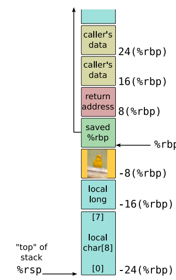
Outline

Canary in the coal mine

Adjacent canary idea



Terminator canary

- Value hard to reproduce because it would tell the copy to stop
- StackGuard: 0x00 0D 0A FF
  - 0: String functions
  - newline: `fgets()`, etc.
  - -1: `getc()`
  - carriage return: similar to newline?
- Doesn't stop: `memcpy`, custom loops

Random canary

- Can't reproduce because attacker can't guess
- For efficiency, usually one per execution
- Ineffective if disclosed

XOR canary

- Want to protect against non-sequential overwrites
- XOR return address with value c at entry
- XOR again with c before return
- Standard choice for c: see random canary

Further refinements

- More flexible to do earlier in compiler
- Rearrange buffers after other variables
  - Reduce chance of non-control overwrite
- Skip canaries for functions with only small variables
  - Who has an overflow bug in an 8-byte array?

What's usually not protected?

- Backwards overflows
- Function pointers
- Adjacent structure fields
- Adjacent static data objects

Where to keep canary value

- Fast to access
- Buggy code/attacker can't read or write
- Linux/x86: `%gs:0x14`

Complex anti-canary attack

- Canary not updated on `fork` in server
- Attacker controls number of bytes overwritten

Complex anti-canary attack

- Canary not updated on `fork` in server
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- ANRY BNRY CNRY DNRY ENRY FNRY
- search $2^{32} \to$ search $4 \cdot 2^8$

Shadow return stack

- Suppose you have a safe place to store the canary
- Why not just store the return address there?
- Needs to be a separate stack
- Ultimate return address protection

Outline

Return address protections

Announcements intermission

ASLR and counterattacks

Testing and fuzzing

Note to early readers

- This is the section of the slides most likely to change in the final version
- If class has already happened, make sure you have the latest slides for announcements

Brief announcements

- Problem set 1 is available on the public web page now
  - Due Friday, 2/17
- The first midterm exam will be next Tuesday (2/21) in class
  - Open book, open notes, no electronics
  - You will have the whole class period
  - Topics will be memory safety bugs and attacks, and threat modeling
  - Similar concepts, but less depth, than labs and p-set

Outline

- Return address protections
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- ASLR and counterattacks
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Basic idea

- "Address Space Layout Randomization"
- Move memory areas around randomly so attackers can't predict addresses
- Keep internal structure unchanged
  - E.g., whole stack moves together

Code and data locations

- Execution of code depends on memory location
- E.g., on x86-64:
  - Direct jumps are relative
  - Function pointers are absolute
  - Data can be relative (`%rip`-based addressing)

Relocation (Windows)

- Extension of technique already used in compilation
- Keep table of absolute addresses, instructions on how to update
- Disadvantage: code modifications take time on load, prevent sharing

PIC/PIE (GNU/Linux)

- "Position-Independent Code / Executable"
- Keep code unchanged, use register to point to data area
- Disadvantage: code complexity, register pressure hurt performance (especially 32-bit)

What's not covered

- Main executable (Linux PIC)
- Incompatible DLLs (Windows)
- Relative locations within a module/area

Entropy limitations

- Intuitively, *entropy* measures amount of randomness, in bits
- Random 32-bit int: 32 bits of entropy
- ASLR page aligned, so at most $32 - 12 = 20$ bits of entropy on x86-32
- Other constraints further reduce possibilities

Leakage limitations

- If an attacker learns the randomized base address, can reconstruct other locations
- Any stack address $\rightarrow$ stack unprotected, etc.

Outline

- Return address protections
- Announcements intermission
- ASLR and counterattacks
- Testing and fuzzing

Testing and security

- "Testing shows the presence, not the absence of bugs" – Dijkstra
- Easy versions of some bugs can be found by targeted tests:
  - Buffer overflows: long strings
  - Integer overflows: large numbers
  - Format string vulnerabilities: `%x`

Random or fuzz testing

- Random testing can also sometimes reveal bugs
- Original 'fuzz' (Miller): `program </dev/urandom`
- Even this was surprisingly effective

Mutational fuzzing

- Instead of totally random inputs, make small random changes to normal inputs
- Changes are called *mutations*
- Benign starting inputs are called *seeds*
- Good seeds help in exercising interesting/deep behavior

Grammar-based fuzzing

- Observation: it helps to know what correct inputs look like
- Grammar specifies legal patterns, run backwards with random choices to generate
- Generated inputs can again be basis for mutation
- Most commonly used for standard input formats
  - Network protocols, JavaScript, etc.

What if you don't have a grammar?

- Input format may be unknown, or buggy and limited
- Writing a grammar may be too much manual work
- Can the structure of interesting inputs be figured out automatically?

Coverage-driven fuzzing

- Instrument code to record what code is executed
- An input is interesting if it executes code that was not executed before
- Only interesting inputs are used as basis for future mutation

AFL

- Best known open-source tool, pioneered coverage-driven fuzzing
- American Fuzzy Lop, a breed of rabbits
- Stores coverage information in a compact hash table
- Compiler-based or binary-level instrumentation
- Has a number of other optimizations