#### CSci 427IW Development of Secure Software Systems Day 25: Crypto failure, authentication

Stephen McCamant University of Minnesota, Computer Science & Engineering

#### Outline

#### SSL/TLS, cont'd

More causes of crypto failure

Announcements intermission

DNSSEC

User authentication

## HTTPS hierarchical PKI

Browser has order of 100 root certs Not same set in every browser

- Standards for selection not always clear
- Many of these in turn have sub-CAs
- 🖲 Also, "wildcard" certs for individual domains

## CA validation standards

CA's job to check if the buyer really is foo.com

#### Race to the bottom problem:

- CA has minimal liability for bad certs
- Many people want cheap certs
- Cost of validation cuts out of profit

#### "Extended validation" (green bar) certs attempt to fix

## HTTPS and usability

Many HTTPS security challenges tied with user decisions

#### Is this really my bank?

Seems to be a quite tricky problem

- Security warnings often ignored, etc.
- We'll return to this as an example later

#### Outline

#### SSL/TLS, cont'd

More causes of crypto failure

Announcements intermission

DNSSEC

User authentication







# Debian/OpenSSL RNG failure (2)

- Debian maintainer commented out some lines to fix a Valgrind warning
  - "Potential use of uninitialized value"
- Accidentally disabled most entropy (all but 16 bits)
- Brief mailing list discussion didn't lead to understanding
- Broken library used for ~2 years before discovery

# Detected RSA/DSA collisions New © 2012: around 1% of the SSL keys on the public net are breakable An Infin © Some sites share complete keypairs An Infin © RSA keys with one prime in common (detected by large-scale GCD) Smaller © One likely culprit: insufficient entropy in key generation Smaller © Embedded devices, Linux /dev/urandom vs. /dev/random Major p Coppers E.g.,

# Newer factoring problem (CCS'17)

- **a** An Infineon RSA library used primes of the form  $p = k \cdot M + (65537^{\alpha} \text{ mod } M)$
- Smaller problems: fingerprintable, less entropy
- Major problem: can factor with a variant of Coppersmith's algoritm
  - E.g., 3 CPU months for a 1024-bit key

# Side-channel attacks

🖲 Timing analysis:

- Number of 1 bits in modular exponentiation
- Unpadding, MAC checking, error handling
- Probe cache state of AES table entries
- 🖲 Power analysis
  - Especially useful against smartcards
- Fault injection
- 🖲 Data non-erasure
  - Hard disks, "cold boot" on RAM

# WEP "privacy"

- First WiFi encryption standard: Wired Equivalent Privacy (WEP)
- F&S: designed by a committee that contained no cryptographers
- Problem 1: note "privacy": what about integrity?
  - Nope: stream cipher + CRC = easy bit flipping



## WEP key size and IV size

- Original sizes: 40-bit shared key (export restrictions) plus 24-bit IV = 64-bit RC4 key
  Both too small
- 🖲 128-bit upgrade kept 24-bit IV
  - Vague about how to choose IVs
  - Least bad: sequential, collision takes hours
  - Worse: random or everyone starts at zero



## Newer problem with WPA (CCS'17)

- Session key set up in a 4-message handshake
- Key reinstallation attack: replay #3
  - Causes most implementations to reset nonce and replay counter
  - In turn allowing many other attacks
  - One especially bad case: reset key to 0
- Protocol state machine behavior poorly described in spec
  - Outside the scope of previous security proofs

## Trustworthiness of primitives

- Classic worry: DES S-boxes
- Obviously in trouble if cipher chosen by your adversary
- In a public spec, most worrying are unexplained elements
- Best practice: choose constants from well-known math, like digits of  $\pi$

## Dual\_EC\_DRBG (1)

- Pseudorandom generator in NIST standard, based on elliptic curve
- Looks like provable (slow enough!) but strangely no proof
- Specification includes long unexplained constants
- Academic researchers find:
  - Some EC parts look good
  - But outputs are statistically distinguishable

## Dual\_EC\_DRBG (2)

- Found 2007: special choice of constants allows prediction attacks
  - Big red flag for paranoid academics
- Significant adoption in products sold to US govt. FIPS-140 standards
  - Semi-plausible rationale from RSA (EMC)
- NSA scenario basically confirmed by Snowden leaks NIST and RSA immediately recommend withdrawal

#### Outline

#### SSL/TLS, cont'd

More causes of crypto failure

Announcements intermission

DNSSEC

User authentication

## Note to early readers

- This is the section of the slides most likely to change in the final version
- If class has already happened, make sure you have the latest slides for announcements

#### Outline

SSL/TLS, cont'd More causes of crypto failure Announcements intermission

DNSSEC

User authentication



## DNSSEC goals and non-goals

- + Authenticity of positive replies
- + Authenticity of negative replies
- + Integrity
- Confidentiality
- Availability



Each resource record gets an RRSIG signature
 E.g., A record for one name-address mapping
 Observe: signature often larger than data

- Signature validation keys in DNSKEY RRs
- Recursive chain up to the root (or other "anchor")



- DNS needs to scale to very large flat domains like . com
- Facilitated by having single DS RR in parent indicating delegation
- Chain to root now includes DSes as well

#### Negative answers

Also don't want attackers to spoof non-existence Gratuitous denial of service, force fallback, etc.

But don't want to sign "x does not exist" for all x

Solution 1, NSEC: "there is no name between acacia and baobab"

## Preventing zone enumeration

- Many domains would not like people enumerating all their entries
- DNS is public, but "not that public"
- Unfortunately NSEC makes this trivial
- Compromise: NSEC3 uses password-like salt and repeated hash, allows opt-out

#### DANE: linking TLS to DNSSEC

"DNS-based Authentication of Named Entities"
DNS contains hash of TLS cert, don't need CAs
How is DNSSEC's tree of certs better than TLS's?

## Signing the root

- Political problem: many already distrust US-centered nature of DNS infrastructure
- Practical problem: must be very secure with no single point of failure
- Finally accomplished in 2010
  - Solution involves 'key ceremonies', international committees, smart cards, safe deposit boxes, etc.



## What about privacy?

- Users increasingly want privacy for their DNS queries as well
- Older DNSCurve and DNSCrypt protocols were not standardized
- More recent "DNS over TLS" and "DNS over HTTPS" are RFCs
- DNS over HTTPS in major browsers might have serious centralization effects



# Passwords: love to hate

Many problems for users, sysadmins, researchers
 But familiar and near-zero cost of entry

User-chosen passwords proliferate for low-stakes web site authentication

#### Password entropy

- Model password choice as probabilistic process
- $\bigcirc$  If uniform,  $\log_2 |S|$
- Controls difficulty of guessing attacks
- Hard to estimate for user-chosen passwords Length is an imperfect proxy





- Online: send guesses to server
- Offline: attacker can check guesses internally
- Specialized password lists more effective than literal dictionaries

 $\blacksquare$  Also generation algorithms (s  $\rightarrow$  \$, etc.)

~25% of passwords consistently vulnerable

## Better password hashing

**Output** Generate random salt s, store (s, h(s, p))

- Block pre-computed tables and equality inferences
- Salt must also have enough entropy
- Deliberately expensive hash function
  - AKA password-based key derivation function (PBKDF)
  - Requirement for time and/or space

## Password usability

- User compliance can be a major challenge
   Often caused by unrealistic demands
- Distributed random passwords usually unrealistic
- Password aging: not too frequently
- Never have a fixed default password in a product

## **Backup authentication**

Desire: unassisted recovery from forgotten password

- Fall back to other presumed-authentic channel Email, cell phone
- Harder to forget (but less secret) shared information
   Mother's maiden name, first pet's name
- 🖲 Brittle: ask Sarah Palin or Mat Honan

## Backup auth suggestion: use time

- Need for backup often comes for infrequently-used accounts
- May be acceptable to slow down recovery if it reduces attack risk
  - Account recovery is a hassle anyway
- Time can allow legitimate owner to notice malicious request

# Centralized authentication

- 🖲 Enterprise-wide (e.g., UMN ID)
- 🗐 Anderson: Microsoft Passport
- 🖲 Today: Facebook Connect, Google ID
- May or may not be single-sign-on (SSO)

# **Biometric authentication**

- Authenticate by a physical body attribute
- + Hard to lose
- Hard to reset
- Inherently statistical
- Variation among people

## **Example biometrics**

- 🖲 (Handwritten) signatures
- 🖲 Fingerprints, hand geometry
- Face and voice recognition
- 🖲 Iris codes