

1. Read the paper Gibbs, N E., Poole, W G., JR., and Stockmeyer, P.K. "An algorithm for reducing the bandwidth and profile of a sparse matrix". SIAM J. Numer. Analysis 13, 2 (April 1976), 235-251. [Posted in the homework page of the class web-site – password required.]

Then in no more than 1/2 page discuss the main differences between the method described in this paper and that used in the standard reverse Cuthill Mc Kee ordering.

2. Consider the matrix (and adjacency graph) shown at the top of Page 6-18 of the lecture notes. 1) Let $\pi = [1, 6, 5, 3, 2, 4]$. Show the adjacency graph and new pattern of the matrix B obtained by symmetrically permuting the original matrix according to the permutation π . 2) Find all fill-ins generated by Gaussian elimination on the matrix B by showing all fill-paths in the graph. 3) Find the Cuthill McKee (CMK) ordering for B (Start at node 1. For tie breaking assume that adjacency lists are sorted increasingly). 4) Find its reverse Cuthill McKee ordering and show the adjacency graph as well as the corresponding reordered matrix C . 5) Repeat question 2 when B is replaced by C .
3. Problem 3-11 from textbook (p. 103 of on-line edition, p. 100 of the SIAM edition). [Hint for question b: 1) You can use induction. 2) A graph that is 2-colorable cannot have a cycle of odd length.]
4. (Question in P. 7-24 of lecture notes). Show that the size of the independent set I found by the greedy algorithm (P. 7-24 of notes) satisfies: $|I| \geq n/(d_I + 1)$, where d_I is the max. degree of all nodes in I (self-cycles (i, i) are not counted). How does this compare with the actual size obtained for a graph associated with a 5-point (rectangular) finite difference grid.
5. * [MATLAB] Recall how we exploited Depth-First Search (DFS) for solving triangular systems with sparse right-hand sides (RHSs). We saw in class that this can be used to implement a sparse LU factorization algorithm based on a left-looking Gaussian elimination (called Gilbert and Peierls algorithm).

The goal of this question is to program the algorithm in matlab by paying attention to efficiency. So the algorithm should be coded in matlab but with an understanding of how you would do it in C/C++, or another programming language.

a. Based on the script `JKI1.m` [posted in matlab page] develop a sparse version of the LU factorization. There are comments in the script indicating the main changes you will need to make. There are two main differences with the dense `JKI` script. First the loop `for k=1:j-1` now becomes `for l=1:length(Lst)` since we are solving only for unknowns in `Lst` and in the order given by `Lst`. The list `Lst` is obtained by a topological sort of the graph with a loop over the nonzero entries of the right-hand side. Second, if you program the code in C, you would allocate the j th column of $[L, U]$ (called w in the notes and the `JKI1.m` script). So you need to explicitly use a temporary column w and when this column is computed copy the appropriate parts to L and U from it.

b. Once you have debugged your script try it on a sample Laplacean problem using the `fd3d` code you already used. Use a $20 \times 20 \times 1$ mesh.

c. Learn how to use profiling in matlab (`profile on`) Profile your execution. How much time is spent where? How would you recode the `dfs` part?

6. The centrality of a node in a graph is a measure of a sort of connectivity of this node with other nodes. The question is: can we use this for reordering an SPD matrix so as to reduce fill-in? A common strategy has been use the degrees [which are updated after each elimination]. Do the following for the matrix *can_256* available from the class web-site:
- Compute the centralities of each node (recall these are the diagonal entries of the matrix $M = \text{expm}(\text{full}(A))$ (in matlab).
 - Define the permutation p that corresponds to increasing centralities and permute the matrix A according to p ($B = A(p, p)$ in Matlab)
 - Do an LU factorization of the permuted matrix.
 - Compare the number of nonzero entries you obtain with those of (i) natural ordering (2) rcm and (3) AMD.