#### GRAPH LAPLACEANS AND THEIR APPLICATIONS

- Graph Laplaceans, definitions and basic properties
- Graph partitioning –
- Introduction to clustering
- Graph Embeddings, vertex embeddings. The problem
- Use of Graph Laplaceans, Laplacean Eigenmaps
- Use of similarity graphs: Locally Linear Embeddings
- Explicit dimension reduction method: PCA, LLP, ...

# Graph Laplaceans - Definition

- "Laplace-type" matrices associated with general undirected graphs useful in many applications
  - ightharpoonup Given a graph G = (V, E) define
    - lacksquare A matrix W of weights  $w_{ij}$  for each edge
    - lacksquare Assume  $w_{ij} \geq 0,$ ,  $w_{ii} = 0,$  and  $w_{ij} = w_{ji} \ orall (i,j)$
    - lacksquare The diagonal matrix  $D = diag(d_i)$  with  $d_i = \sum_{j 
      eq i} w_{ij}$
- Corresponding graph Laplacean of G is:

$$L = D - W$$

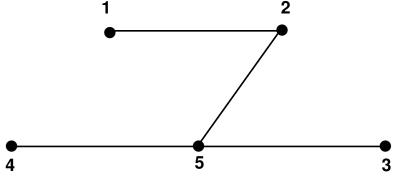
ightharpoonup Gershgorin's theorem ightarrow L is positive semidefinite.

> Simplest case:

$$w_{ij} = egin{cases} 1 & ext{if } (i,j) \in E\&i 
eq j \ 0 & ext{else} \end{cases} \quad D = ext{diag} \left[ d_i = \sum_{j 
eq i} w_{ij} 
ight]$$

#### Example:

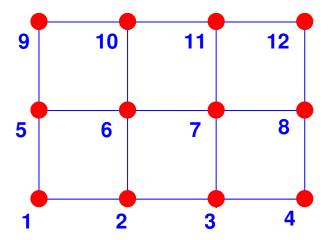
Consider the graph



$$L = \begin{bmatrix} 1 & -1 & 0 & 0 & 0 \\ -1 & 2 & 0 & 0 & -1 \\ 0 & 0 & 1 & 0 & -1 \\ 0 & 0 & 0 & 1 & -1 \\ 0 & -1 & -1 & 3 \end{bmatrix}$$

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Define the graph Laplacean for the graph associated with the simple mesh shown next. [use the simple weights of 0 or 1]. What is the difference with the discretization of the Laplace operator for case when mesh is the same as this graph?



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#### Proposition:

- (i) *L* is symmetric semi-positive definite.
- (ii) L is singular with 1 as a null vector.
- (iii) If G is connected, then  $Null(L) = span\{ 1 \}$
- (iv) If G has k>1 connected components  $G_1,G_2,\cdots,G_k$ , then the nullity of L is k and Null(L) is spanned by the vectors  $z^{(j)},\,j=1,\cdots,k$  defined by:

$$(z^{(j)})_i = \left\{egin{array}{ll} 1 & ext{if } i \in G_j \ 0 & ext{if not.} \end{array}
ight.$$

- Glaplacians

Proof: (i) and (ii) seen earlier and are trivial. (iii) Clearly u=1 is a null vector for L. The vector  $D^{-1/2}u$  is an eigenvector for the matrix  $D^{-1/2}LD^{-1/2}=I-D^{-1/2}WD^{-1/2}$  associated with the smallest eigenvalue. It is also an eigenvector for  $D^{-1/2}WD^{-1/2}$  associated with the largest eigenvalue. By the Perron Frobenius theorem this is a simple eigenvalue... (iv) Can be proved from the fact that L can be written as a direct sum of the Laplacian matrices for  $G_1, \dots, G_k$ .

# A few properties of graph Laplaceans

*Define:* oriented incidence matrix H: (1)First orient the edges  $i \sim j$  into  $i \to j$  or  $j \to i$ . (2) Rows of H indexed by vertices of G. Columns indexed by edges. (3) For each (i,j) in E, define the corresponding column in H as  $\sqrt{w(i,j)}(e_i-e_j)$ .

**Example:** In previous example (4 p. back) orient  $i \rightarrow j$  so that j > i [lower triangular matrix representation]. Then matrix H is:

$$H = egin{bmatrix} 1 & 0 & 0 & 0 \ -1 & 1 & 0 & 0 \ 0 & 0 & 1 & 0 \ 0 & 0 & 0 & 1 \ 0 & -1 & -1 & -1 \end{bmatrix}$$

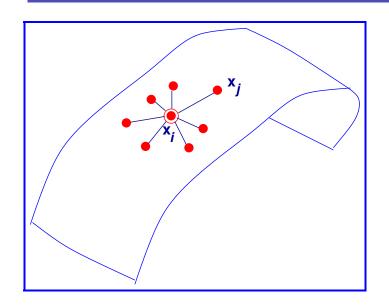
Property 1

$$L = HH^T$$

Re-prove part (iv) of previous proposition by using this property.

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### A few properties of graph Laplaceans



Strong relation between  $x^T L x$  and local distances between entries of x

Let L = any matrix s.t. L = D - W, with  $D = diag(d_i)$  and

$$w_{ij} \geq 0, \qquad d_i \ = \ \sum_{j 
eq i} w_{ij}$$

*Property 2:* for any  $x \in \mathbb{R}^n$ :

$$x^ op L x = rac{1}{2} \sum_{i,j} w_{ij} |x_i - x_j|^2$$

*Property 3:* (generalization) for any  $Y \in \mathbb{R}^{d \times n}$ :

$$\mathsf{Tr}\left[oldsymbol{Y} L oldsymbol{Y}^ op
ight] = rac{1}{2} \sum_{i,j} oldsymbol{w}_{ij} \|oldsymbol{y}_i - oldsymbol{y}_j\|^2$$

Note:  $y_j = j$ -th column of Y. Usually d < n. Each column can represent a data sample.

**Property 4:** For the particular  $L = I - \frac{1}{n} \mathbb{1} \mathbb{1}^{\top}$ 

$$XLX^{\top} = ar{X}ar{X}^{\top} == n imes \mathsf{Covariance}$$
 matrix

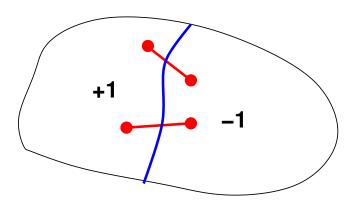
Property 5: L is singular and admits the null vector

$$1 = ones(n, 1)$$

*Property 6:* (Graph partitioning) Consider situation when  $w_{ij} \in \{0,1\}$ . If x is a vector of signs  $(\pm 1)$  then

$$x^ op Lx = 4 imes$$
 ('number of edge cuts') edge-cut = pair  $(i,j)$  with  $x_i 
eq x_j$ 

Consequence: Can be used to partition graphs



- ightharpoonup Would like to minimize (Lx,x) subject to  $x\in\{-1,1\}^n$  and  $e^Tx=0$  [balanced sets]
- ➤ WII solve a relaxed form of this problem

What if we replace x by a vector of ones (representing one partition) and zeros (representing the other)?

Let x be any vector and  $y = x + \alpha$  1 and L a graph Laplacean. Compare (Lx, x) with (Ly, y).

- ightharpoonup Consider any symmetric (real) matrix A with eigenvalues  $\lambda_1 \leq \lambda_2 \leq \cdots \leq \lambda_n$  and eigenvectors  $u_1, \cdots, u_n$
- Recall that: (Min reached for  $x = u_1$ )

$$\min_{x\in\mathbb{R}^n}rac{(Ax,x)}{(x,x)}=\lambda_1$$

In addition: (Min reached for  $x = u_2$ )

$$\min_{x\perp u_1}rac{(Ax,x)}{(x,x)}=\lambda_2$$

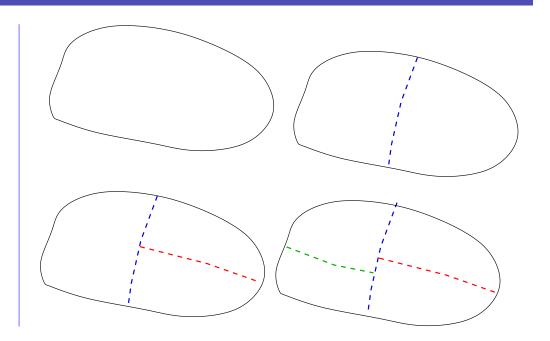
- For a graph Laplacean  $u_1 = 1 = 1$  vector of all ones and
- ightharpoonup ....vector  $u_2$  is called the Fiedler vector. It solves a relaxed form of the problem -

$$\min_{oldsymbol{x} \in \{-1,1\}^n; \ \ extstyle ^T x = 0} rac{(Lx,x)}{(x,x)} 
ightarrow \sum_{oldsymbol{x} \in \mathbb{R}^n; \ \ extstyle ^T x = 0} rac{(Lx,x)}{(x,x)}$$

ightharpoonup Define  $v=u_2$  then lab=sign(v-med(v))

# Recursive Spectral Bisection

- Form graph Laplacean
- Partition graph in 2 based on Fielder vector
- 3 Partition largest subgraph in two recursively ...
- 4 ... Until the desired number of partitions is reached



### Three approaches to graph partitioning:

- 1. Spectral methods Just seen + add Recursive Spectral Bisection.
- 2. Geometric techniques. Coordinates are required. [Houstis & Rice et al., Miller, Vavasis, Teng et al.]
- 3. Graph Theory techniques multilevel,... [use graph, but no coordinates]
  - Currently best known technique is Metis (multi-level algorithm)
  - Simplest idea: Recursive Graph Bisection; Nested dissection (George & Liu, 1980; Liu 1992]
  - Advantages: simplicity no coordinates required

Run testBis\_simple and testMeshPart in matlab class site

# Example of a graph theory approach

- Level Set Expansion Algorithm
- ightharpoonup Given: p nodes 'uniformly' spread in the graph (roughly same distance from one another).
- Method: Perform a level-set traversal (BFS) from each node simultaneously.
- ightharpoonup Best described for an example on a 15 imes 15 five point Finite Difference grid.
- See [Goehring-YS '94, See Cai-YS '95]
- Approach also known under the name 'bubble' algorithm and implemented in some packages [Party, DibaP]

